



Secure Communication without Key Exchange



Honours Project 2013

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Seminar Overview

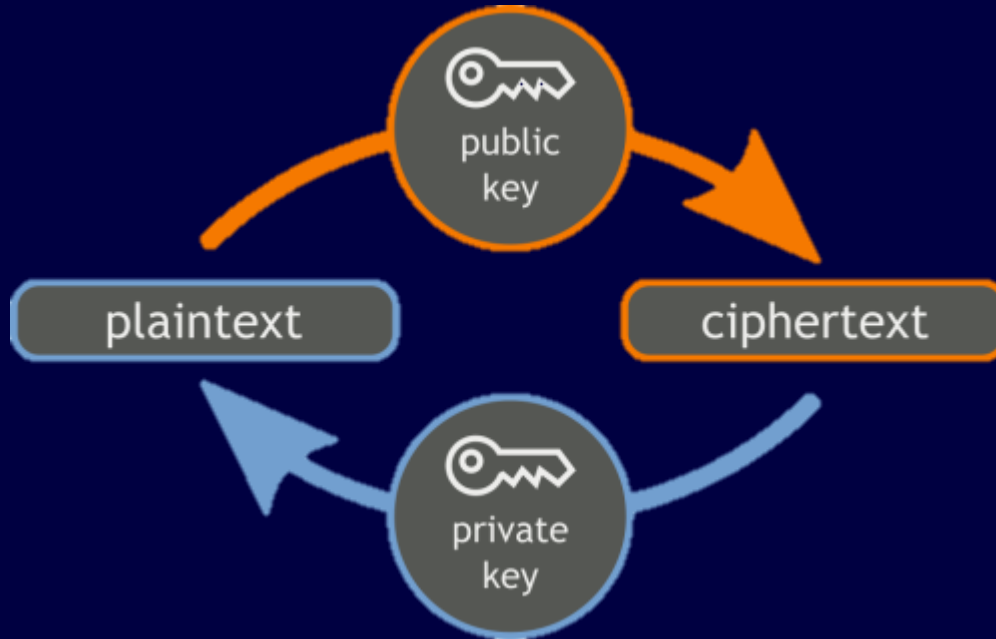
- Objectives & Context
- Project Significance & Implications
- One-Time Pad in the KS Cipher
- Geometric Algebra 3D & 4D
 - Introduction
 - Analytical Work & C++ Program
 - CLUViz Demonstration
 - Summary
- Timing-Based Physical Layer Encryption
- Project Management
- Conclusion
- Question Time & Key References





Introduction

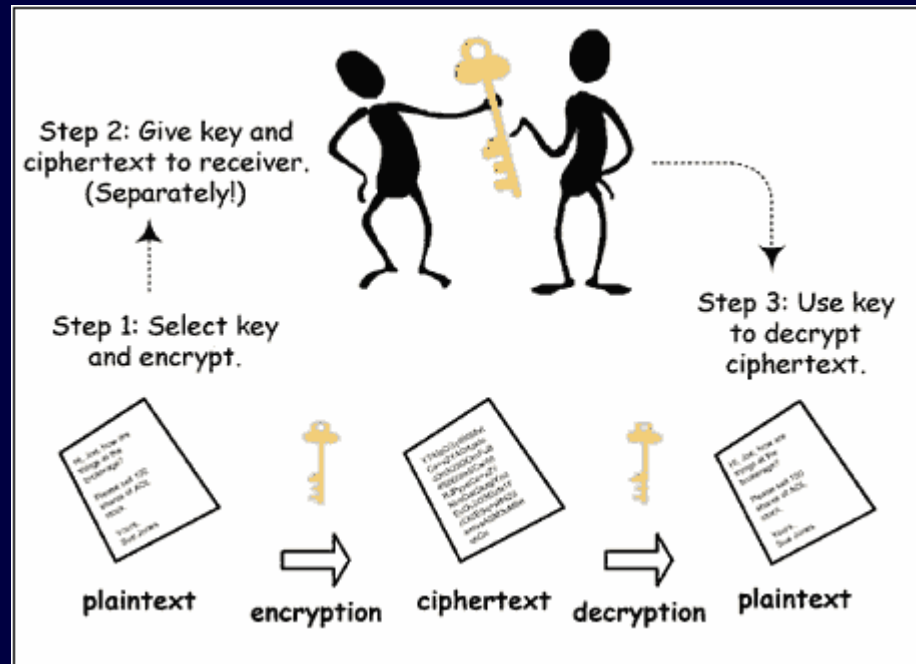
Public Key Cryptography





Introduction

Symmetric Key Systems

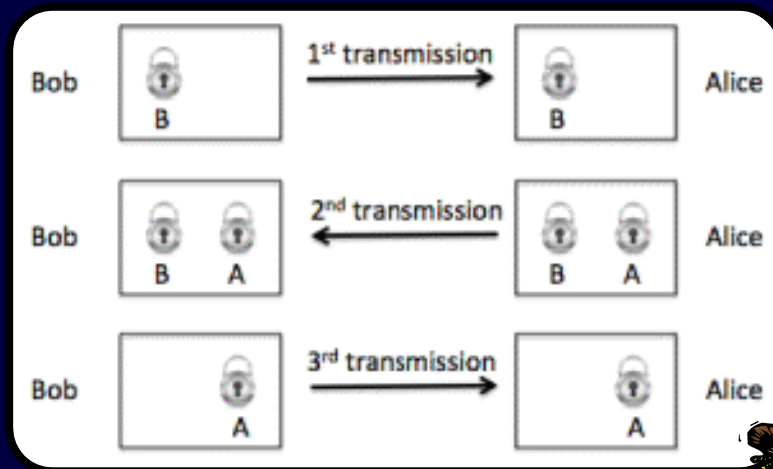




The Double-Padlock Protocol

Kish-Sethuraman (KS) Cipher - The Double Padlock Protocol

What it would mean?



Laszlo Bela Kish
Professor, Texas A&M University

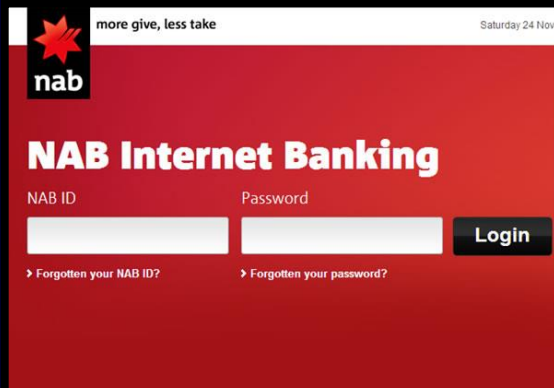


Image: J. Chappell and D. Abbott, The Double -Padlock Problem, <http://scholar.google.com.au/citations?user=fVJ8twEAAA&hl=en>





Project Significance & Implications



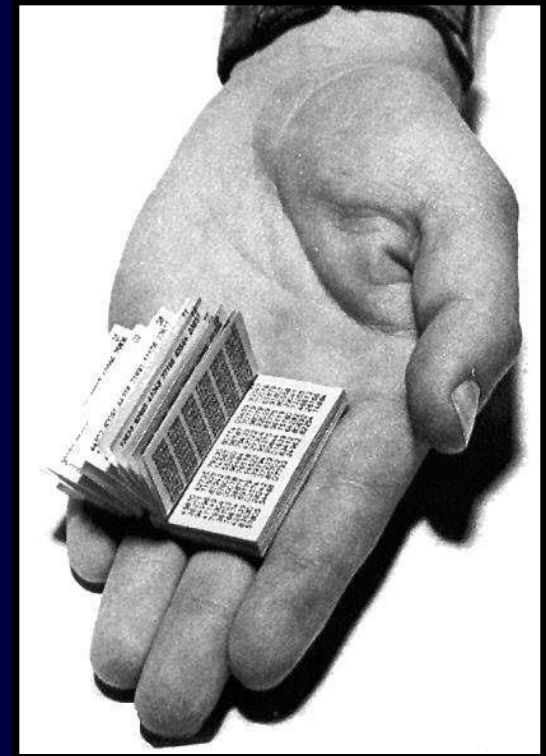
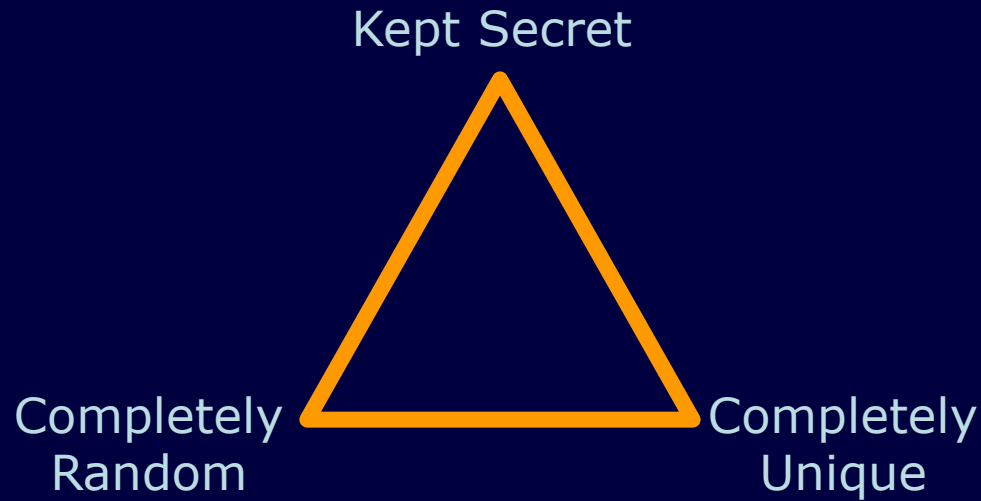
Images: aerospace.firetench.com, <http://www.nab.com.au>, http://www.jvi.org/wp-content/uploads/GPS_Satellite_NASA_art-iif.jpg





The One-Time Pad

Key Elements of the OTP



Small Russian One-Time Pad captured by MI-5

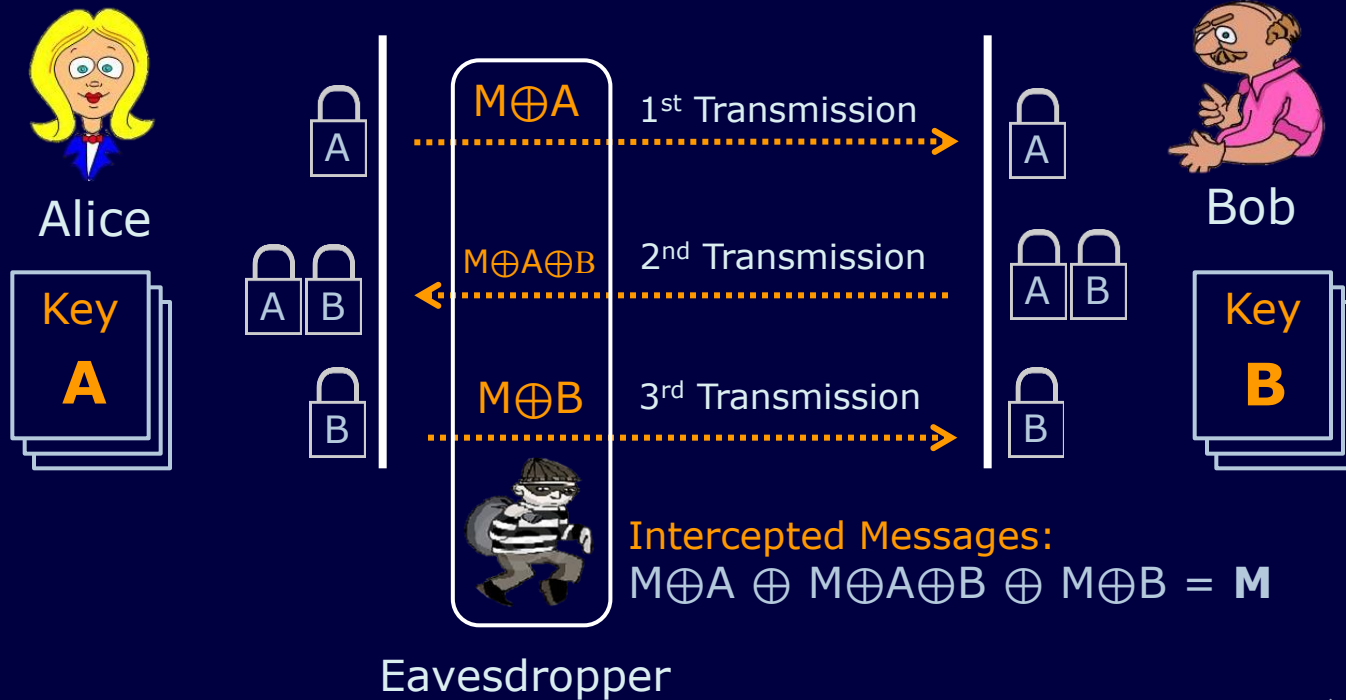
Image: www.ranum.com/security/computer_security/papers/otp-faq/





Using a One Time Pad in the KS-cipher

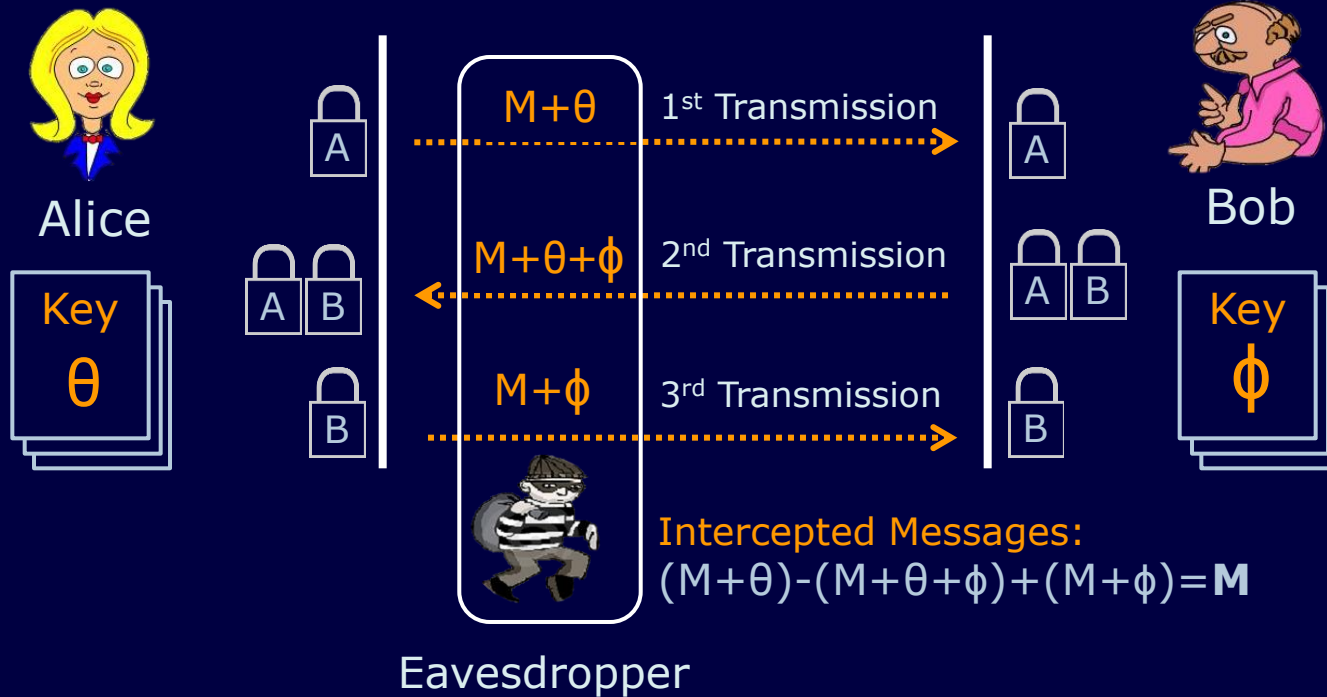
What if Alice & Bob each had their own unique OTP?
The initial Message is M





Using a One Time Pad in the KS-cipher

The XOR approach can be generalised to rotations in 2D

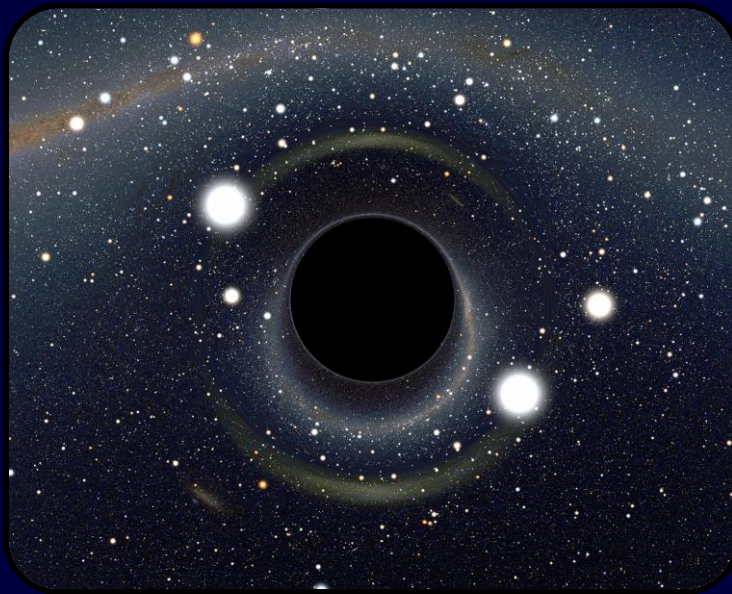




Geometric Algebra

A Powerful Mathematical Tool

Ability to easily handle rotations in N-dimensions



$$v' = e^{\frac{ir\theta}{2}} \cdot v \cdot e^{\frac{-ir\theta}{2}}$$

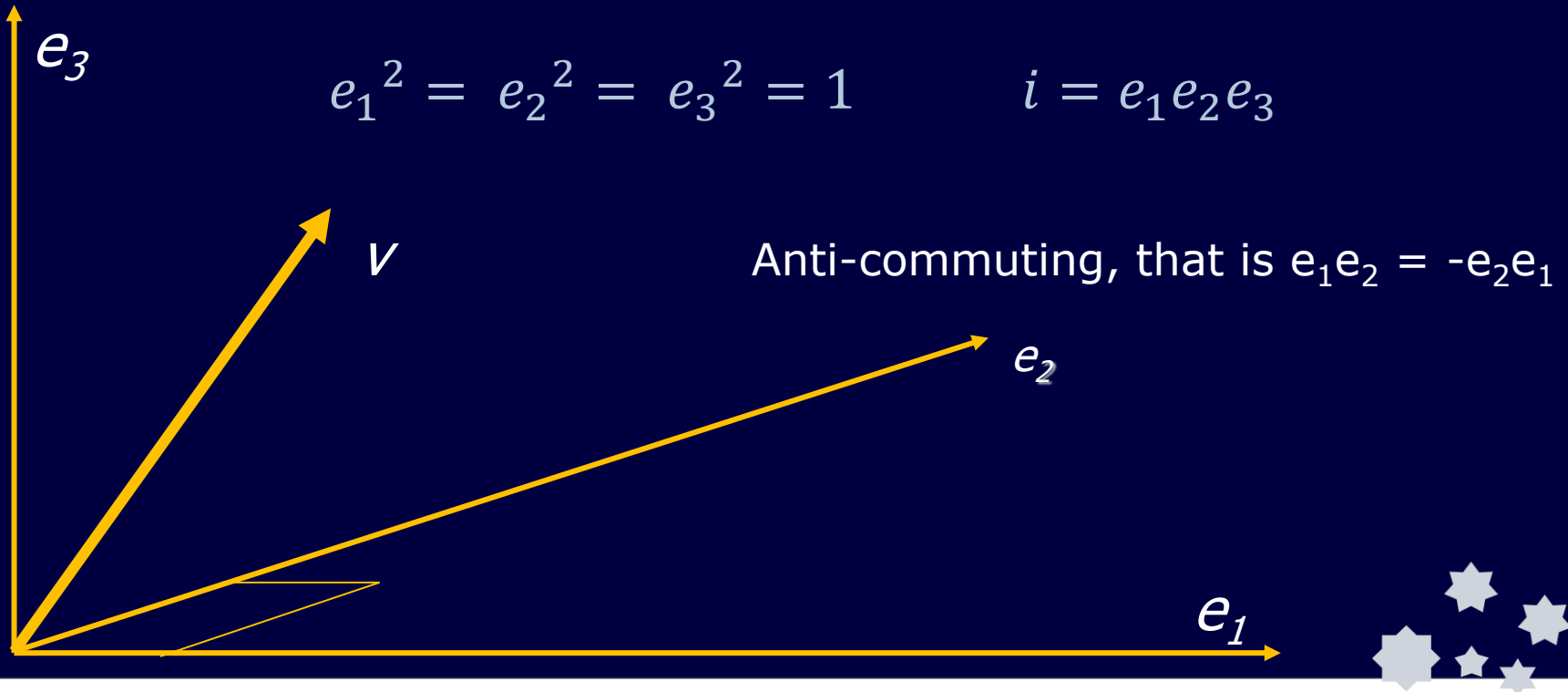
De Moivre's theorem applies





Geometric Algebra

Vector v is defined as $v = a_1e_1 + a_2e_2 + a_3e_3$



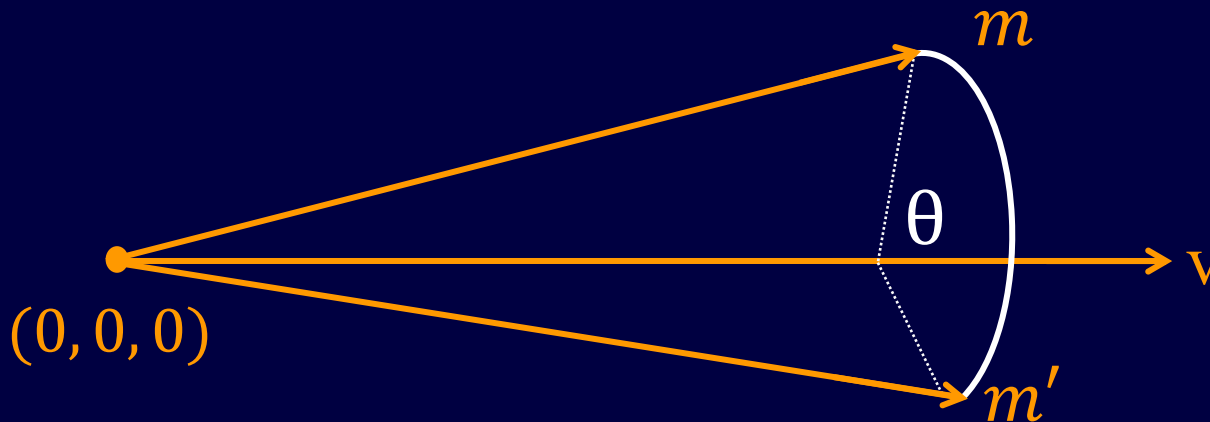


Geometric Algebra

Initial Message Vector $\mathbf{m} = m_1\mathbf{e}_1 + m_2\mathbf{e}_2 + m_3\mathbf{e}_3$

$$R = e^{iv\theta/2}$$

$$\therefore \mathbf{m}' = e^{iv\theta/2} \cdot \mathbf{m} \cdot e^{-iv\theta/2}$$





Geometric Algebra

Rotation Operators of Alice & Bob

$$R_A = e^{iv\theta/2} \quad R_B = e^{iw\phi/2}$$

$$\mathbf{m}_{final} = \widetilde{R}_B \widetilde{R}_A R_B R_A \mathbf{m}_{initial} \widetilde{R}_A \widetilde{R}_B R_A R_B$$

$$R_A R_B - R_B R_A = -\sin\frac{\phi}{2} \sin\frac{\theta}{2} v \times w \quad = 0$$





Geometric Algebra

- C++ Program
- 4D Analytical Work
- Still wanted to explore and visualize 4D Rotations



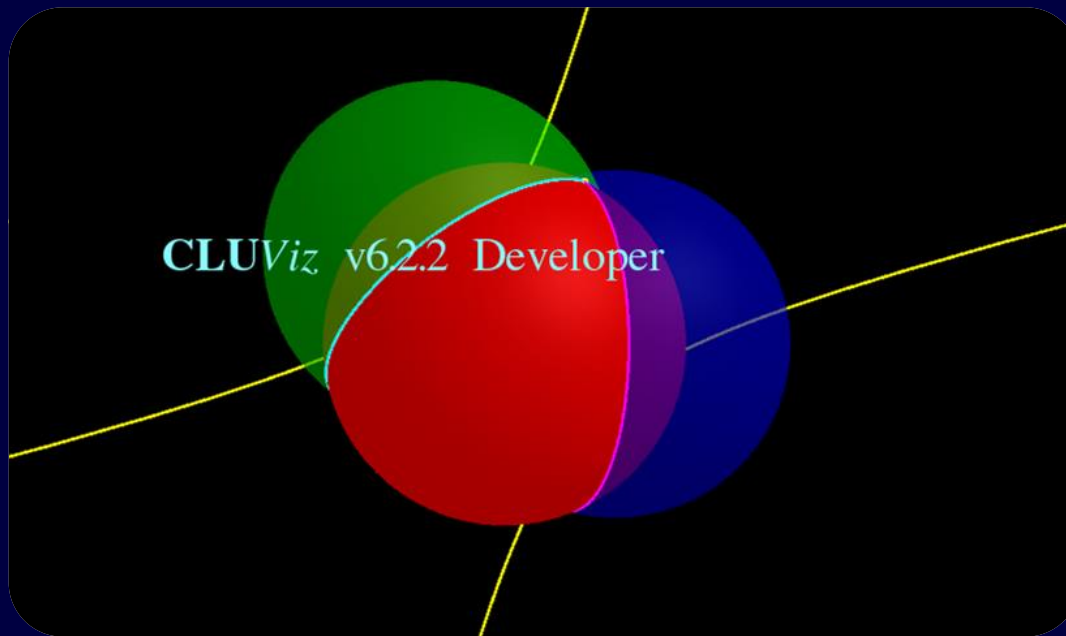
Source: James Chappell





CLUViz Demonstration

- 3D CLUViz Program
- Proof of requirement for parallel rotation axes
 - Complexity of 4D Rotations





Geometric Algebra Summary

Information Theory revealed a hole in our approach

Shannon's work on the capacity of a *Binary Symmetric Channel*

Secrecy Rate $\rightarrow C_S$

Alice's Information $\rightarrow X$

Bob's Information $\rightarrow Y$

$$C_S \leq I(X, Y)$$

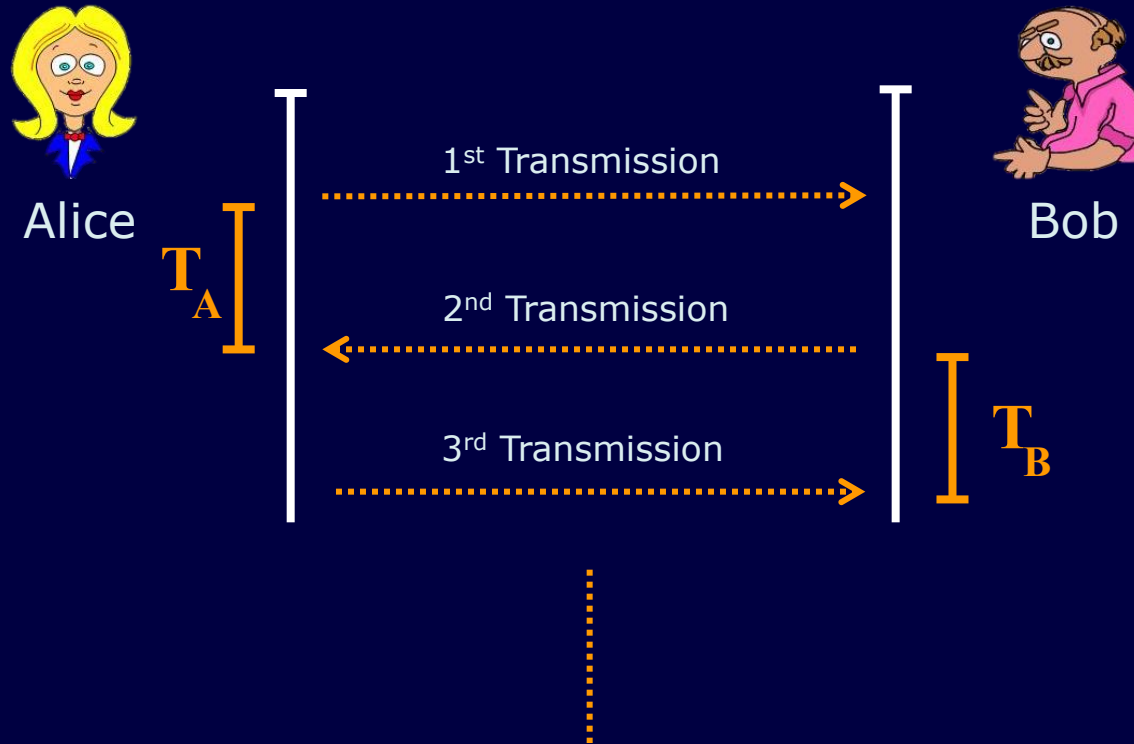
Found a new approach in Timing-based Physical Layer Encryption





Timing-Based Physical Layer Encryption

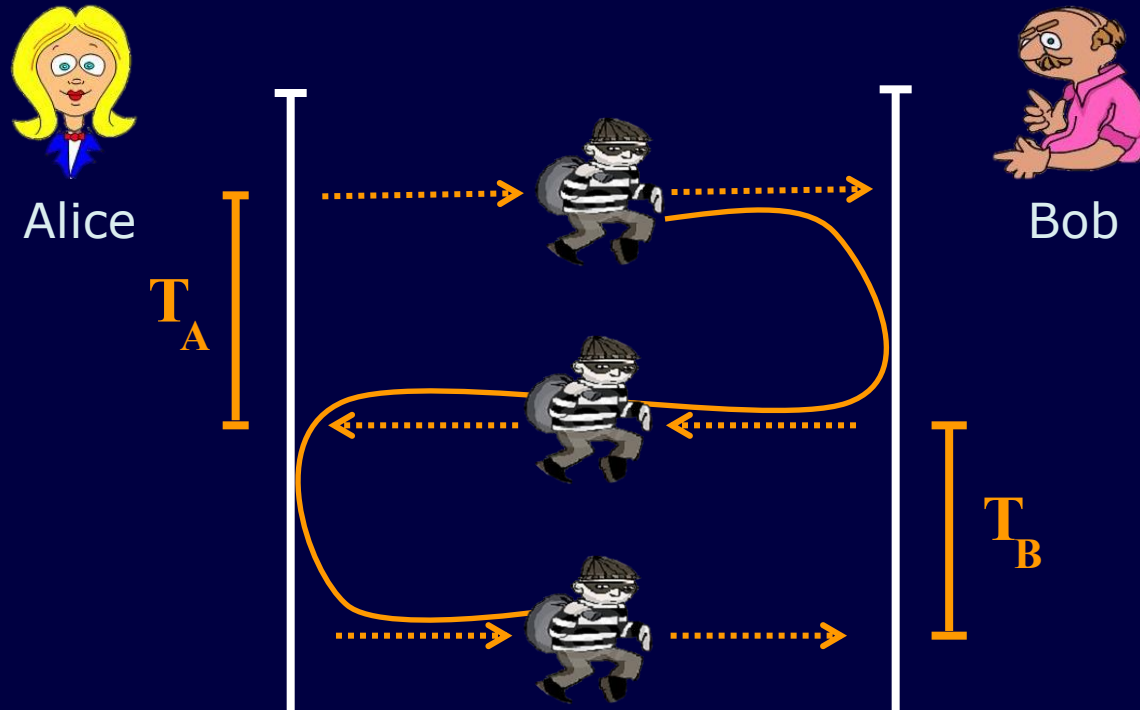
What is it? How can it be used?





Timing-Based Physical Layer Encryption

Practical Setup





Timing-Based Physical Layer Encryption

Bit Stream Generation

$\left(\begin{array}{c} 0.363 \\ 0.407 \\ 0.356 \\ 0.333 \\ 0.565 \\ 0.345 \end{array} \right)$

$\rightarrow 110010\dots$

$rtt(i) > median \rightarrow 1$

$rtt(i) < median \rightarrow 0$





Timing-Based Physical Layer Encryption

Bit Parity Checks

Alice 10 01 11 01 → 1 0 1

Bob 11 01 10 10 → 1 0 1

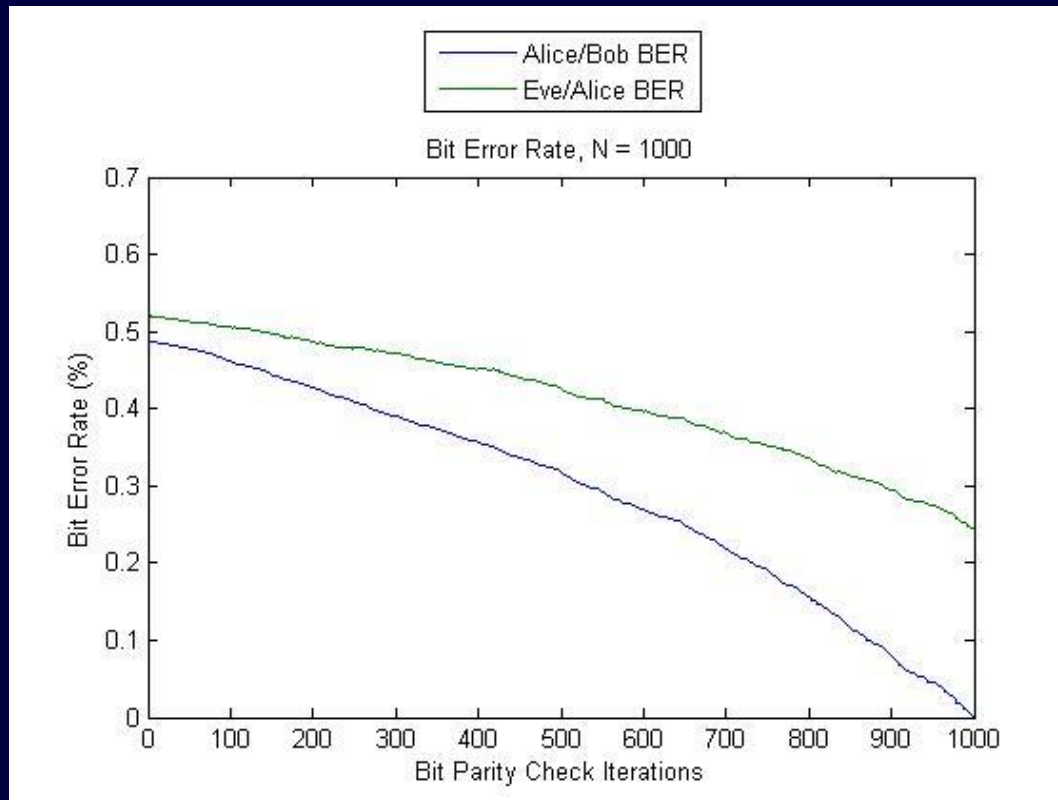
Eve 01 01 00 11 → 0 0 0





Timing-Based Physical Layer Encryption

Bit Parity Checks





Timing-Based Physical Layer Encryption

Matlab Analysis of Output

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Comparison Tool - D:\Uni\Honours\testing1.txt vs. D:\Uni\Honours\testing2.txt
COMPARISON VIEW
+ Refresh Swap Sides Save As Print Find Comparison Help
testing1.txt vs. testing2.txt
37 differences found. Use the tooltip buttons to navigate to them.
1 1001001111110001001101001010001011110100011111100110111010100010001110100101 x 100100111111000100110100101000101111011101111000110111010100010001110100101
2 01101100100111001101000011001000000101100000010011011110000000001100101000 . 01101100100111001101000011001000000101100000010011011110000000001100101000
3 1111101101100100101100110111011001010011100110000001110110110110100100010 . 1111101101100100101100110110111011001010011100110000001110110110110100100010
4 0010000111001001100010001000010001100011111111101011001110011000100011000 x 00100001110010011000101010000100110001100011001111111110101110011000110001100
5 1100001011100111010010010101000010101001010011011110000010110001010010101 . 1100001011100111010010010101000010101001010011011110000010110001010010101
6 1100110000111110110100110100110111100001010000011000010101010010010010101000 . 11001100001111101101001110100110101111000010110000011000010101010010010101000
7 10110110001100100111101101011101101011101010101001010111001101101010011 x 101101100011001001111011010111010101110101011101010110100110011100110101010011
8 11010100110010000000111101001110101011110100111111101001000010000010001110001 x 1101010011001000000111110100111101010111101001111111010010000100001100011
9 11101011111010101101101100100011110101110000111110010100101010001010011100011 . 11101011111010101010101001000111101011100001111100101001010001010011100011
10 10001010111101001000000101011011011010011010001110100001001111000010100110010 . 100010101111010010000001010110110110100110110001110100001001111000010100110010
11 010000100000110001000101110010000001100001100000000100010000000001111010000011100 x 01000010000011000100010111001000000110000110000000010011000000001111010000011100
12 111000110110100001110101100111001110011110110100001001010111010010100010101110100 . 11100011011010000111010110011100111101101000010010101110100101000101011110100
13 0011011001001101001000100010010000011011101101101110101110100100000010001 x 0011011001001101001000100010010000010011011101101110101110101110100000010001
14 0010111100010101000111000001110011110000000001011110001001010001110011100111010 . 00101111000101010001110000011100111100000000010111100010010110001110011101010
15 0010011110011101111001110100001001110100001001010110100001110010000001100100 . 00100111100111011110011101000010011101000010010101010000111001000001100100
16 11100100111111100111100100010000001101001100111101101101100101011000110011 . 1110010011111110011110010001000000110100110011110110110110100010110100011001
17 01010011100010110110001000001110001100001000101110001011101000010100010010 . 0101001110001011011000100000111000110000100010101100001011101000010100010010
18 011110000100101000111010100000011010110110010000000001000011101101010010111011 x 01111000011100010001001101001101111101100100000000100001110110101010010111011
19 0010101101011100000101000100110111110100101000100101010100111000001000 x 001010110111100000101000100110101111111010011000100101010100111000001000
20 0010010111000101100110111101011010101001000011110001010000011000010011110110010 x 000001011100010110011011110101101010100100010000111110001101000011000010011110110010
21 11101100101100111011000000000110011000001101111000010100010111101000110010111011 . 1110110010110011101100000000011001100000110111000010100010111101000110010111011
22 1010011011110110101101000000100001111011111111100110100111001001100001111101010 . 1010011011110110101101000000100001111011111111100110100111001001100001111101010

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The length of the two bit streams is 14324
The number of errors between the two bit streams is 195
The BER between Bob and Alice is 1.361%
fx >>

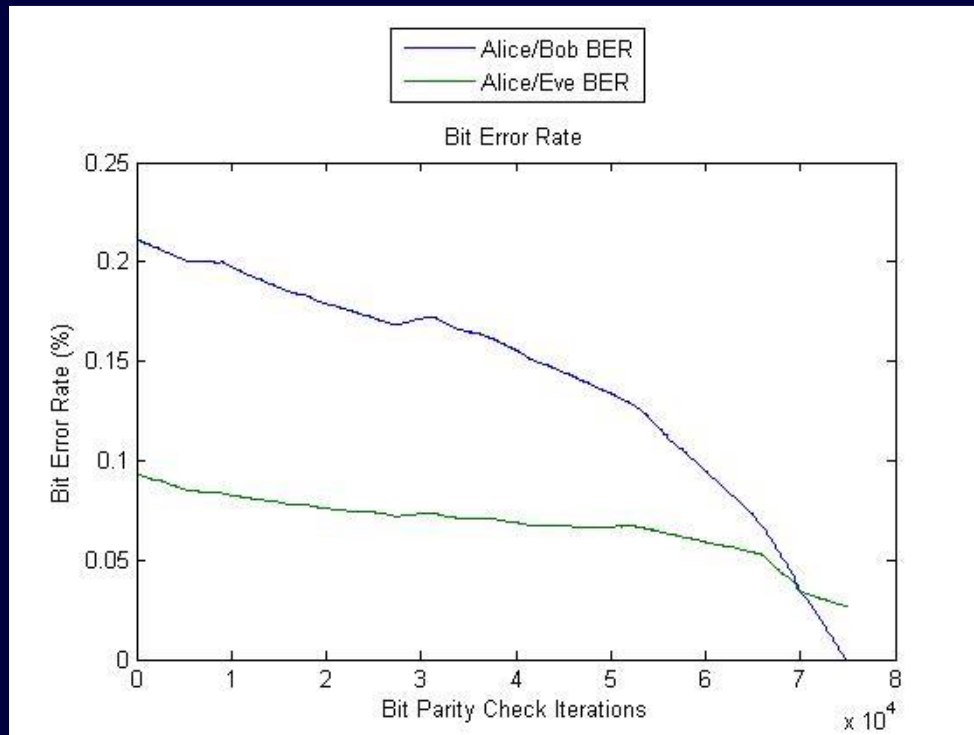
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Timing-Based Physical Layer Encryption

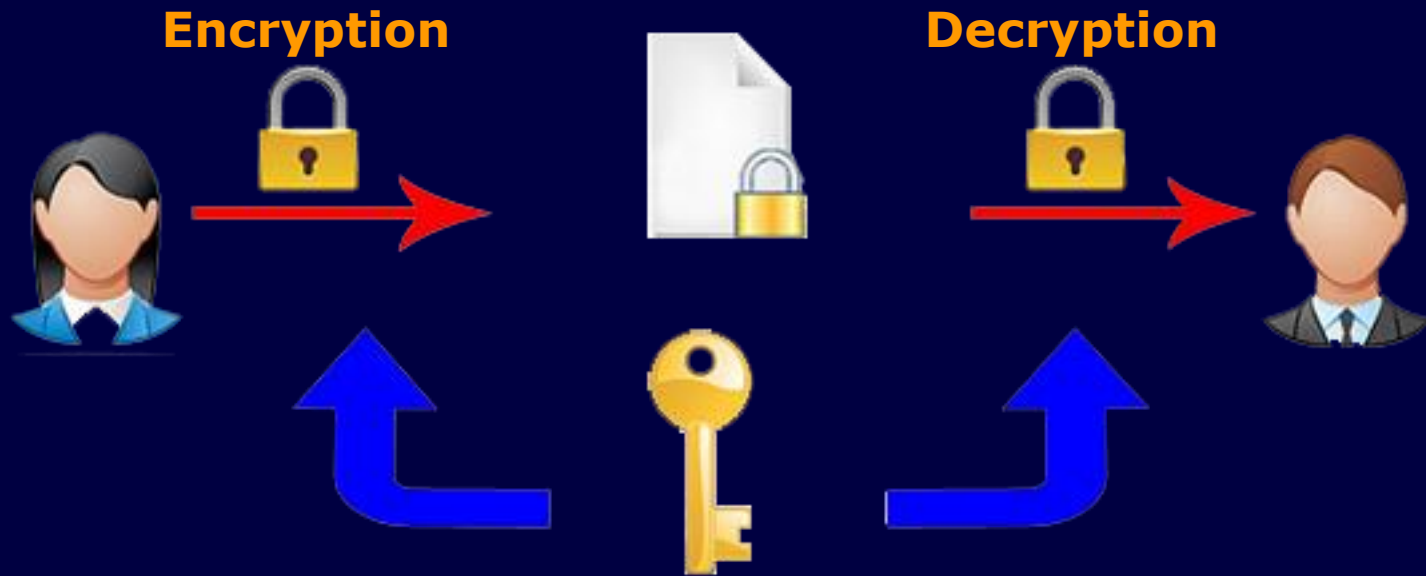
Matlab BER Results





Timing-Based Physical Layer Encryption

Application





Project Management

Risk Management Evaluation

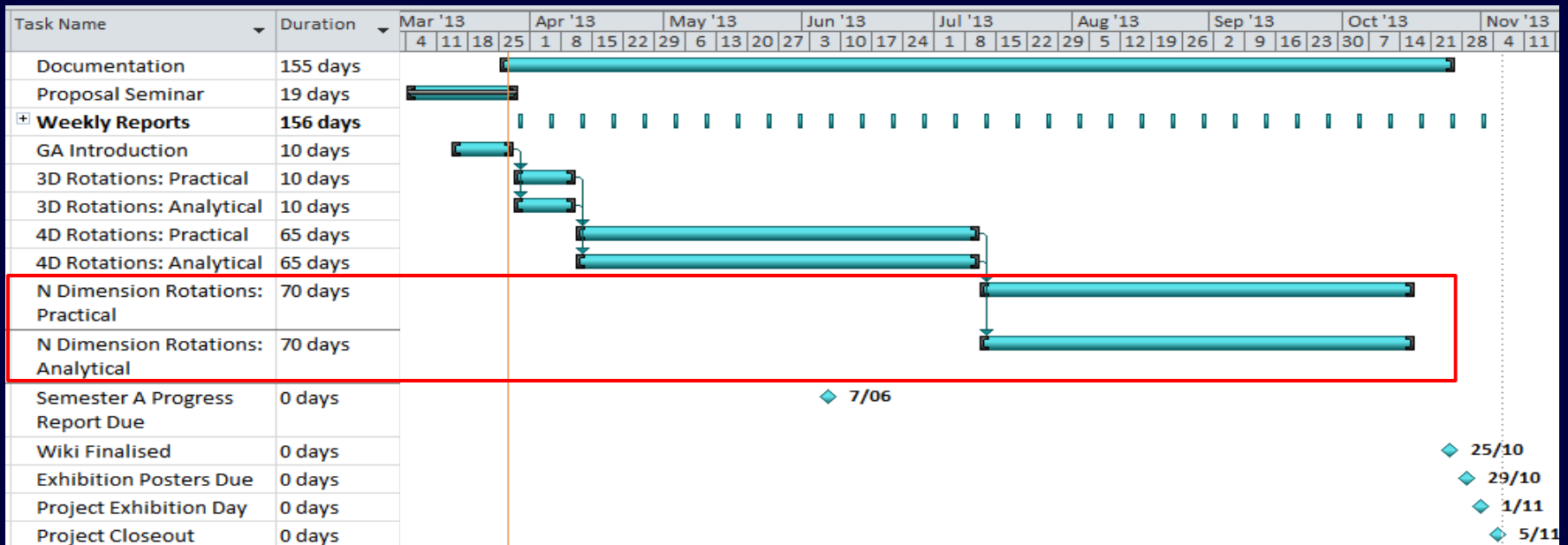
Risk	Likelihood	Severity	Avoidance/Mitigation Strategies
Unavailability of Team Member	Low	Medium	Both members are well versed in each aspect of the project and the overall progress.
Loss of work	Very Low	Low	Group members will ensure that all progress is shared via the wiki, Facebook and through email so that we have several working copies available.
A Lack of Technical Knowledge	Low	High	We'll need to ensure that we're maintaining communication with each other and our supervisors to make sure that we understand the technical elements of the project
Falling Behind Schedule as a result of the increased complexity of the project.	Low	Medium	Re-evaluate our expectations of the project, and perhaps increase the focus in lower dimensions (such as 4, 5 and 6) before even considering the higher dimensions.
Not finding a solution for keyless encryption	Very High	Very Low	Ensure that our work is completely documented, so that regardless of what we've found we have something to show at the project closing.





Project Management

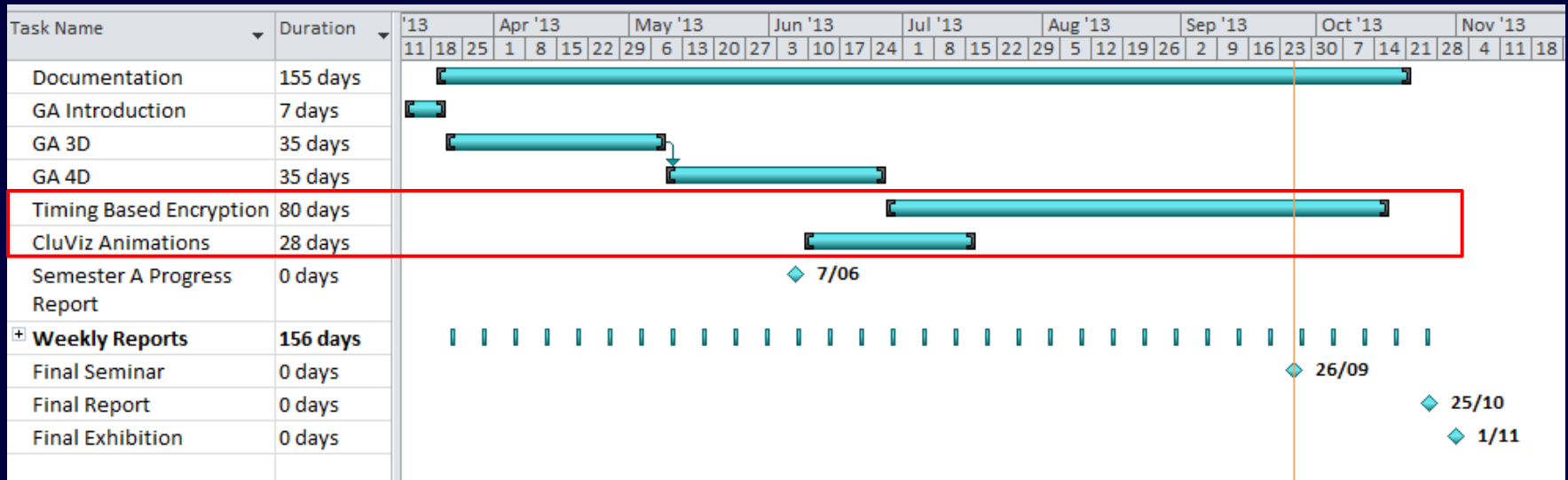
Initial Project Schedule





Project Management

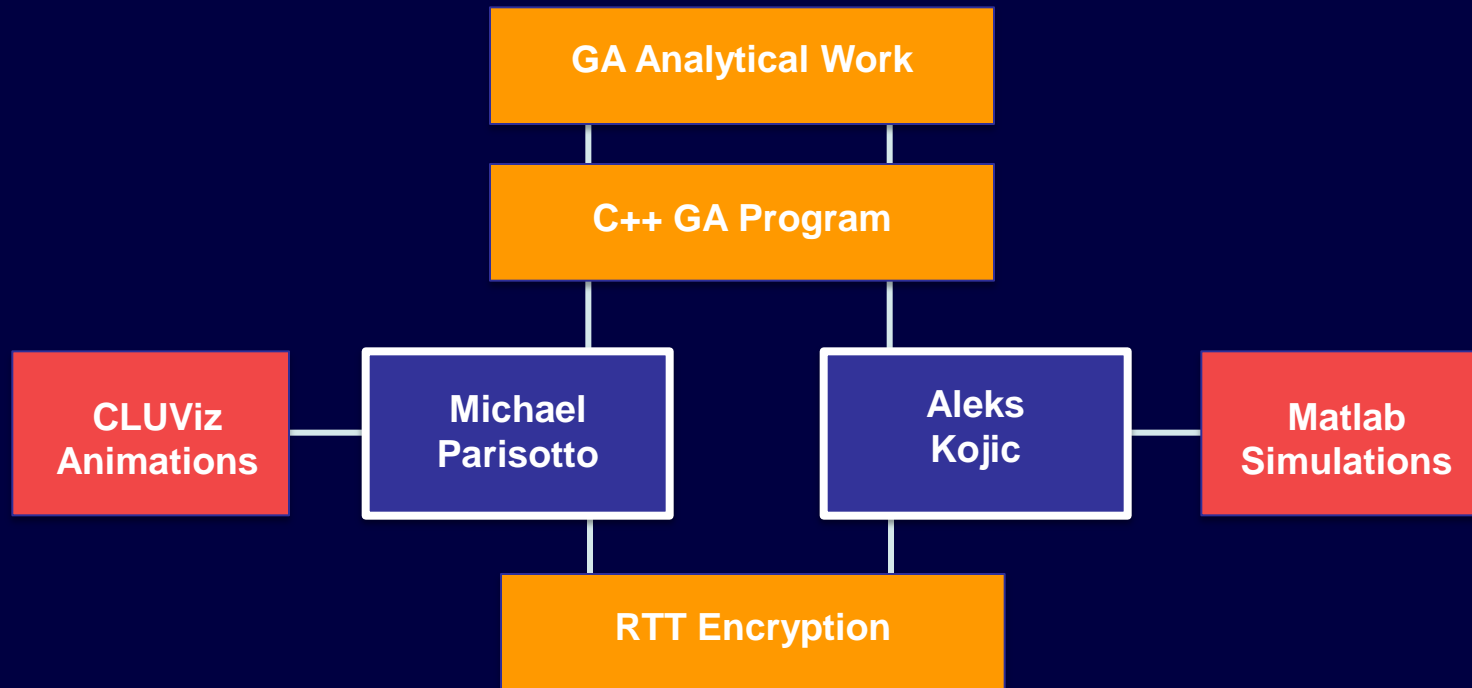
Resulting Project Schedule





Project Management

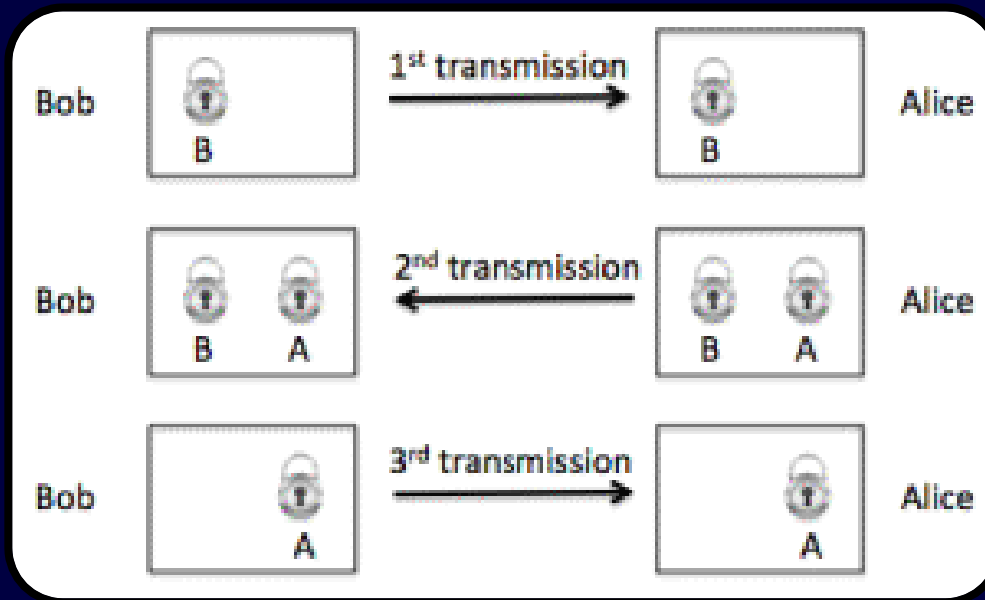
Team Management & Organisation





Conclusion

- Geometric Algebra and Physical Timing Based Encryption
 - Appreciation for Cryptography



Thank you for listening





Questions

Key References

- [1] L. Kish and S. Sethuraman, 'Non-breakable Data Encryption with Classical Information', http://ee.tamu.edu/~noise/research_files/new_encryption.pdf
- [2] L. Kish and J. Bergou, 'Absolutely secure QKD Scheme with no detection noise, entanglement and classical communication', <http://arxiv.org/pdf/quant-ph/0509097.pdf>
- [3] L. Gunn 'Physical-layer encryption on the public internet: a stochastic approach to the KS cipher', <http://arxiv.org/pdf/1306.4174v1.pdf>
- [4] M. Gander and U. Maurer 'The secret-key rate of binary random variables', <http://ieeexplore.ieee.org/stamp/stamp.jsp?tp=&arnumber=394667>

